

**AMENDMENTS TO THE SPECIFICATION:**

Please replace the paragraph beginning on page 3, line 3 with the following amended paragraph:

However, although the conventional RSA encryption scheme as described above has a high performance in terms of data secrecy and has a simple algorithm ~~algorism~~, its security depends on the difficulty of factoring a product  $n$  of two prime numbers  $p$  and  $q$ . Therefore, it is necessary to use about 200-digit  $n$  in the decimal system, and there is the problem that it is very difficult to perform modulo  $n$  exponentiation, which are necessary for encryption and decryption processes.

Please replace the paragraph beginning on page 3, line 17 with the following amended paragraph:

An object of the present invention is to suggest a secret cryptosystem of an extremely simple public key, which simplifies its algorithm ~~algorism~~, while maintaining a security equivalent to the RSA encryption scheme, and to provide an encrypting device which can perform encryption by simple calculations, a decrypting device which can perform decryption by simple calculations, a cryptosystem including the same devices, an encrypting method, and a decrypting method.

Please replace the paragraph beginning on page 20, line 6 with the following amended paragraph:

The cryptosystem of the present embodiment, which has the same basic principle as that of the cryptosystem in the Embodiment 1, can make its algorithm ~~algorism~~ simpler under the condition that a size  $b$  of a private key  $p$  is limited in relation to a message  $m$ .

Please replace the paragraph beginning on page 43, line 15 with the following amended paragraph:

According to the above arrangement, in the encrypting ~~encryptin~~ device, keys  $g_1$  and  $g_2$  generated as a public key respectively include the power of  $(p-1)$  and the power of  $(q-1)$ , and the ciphertext elements  $C_1$  and  $C_2$  generated using the public key  $\{g_1, g_2\}$  and the private key  $n$  also include the power of  $(p-1)$  and the power of  $(q-1)$ , respectively. This makes it possible for the decrypting device to easily decrypt the ciphertext elements  $C_1$  and  $C_2$  using the Fermat's little theorem ( $a^{p-1} \equiv 1 \pmod{p}$ ).